

# Gossip and Knowledge

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▶ Dynamic Gossip

presentation for ESSLLI Galway

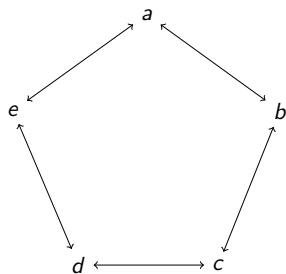
August 2022

## Gossip Graphs — dynamic gossip

We assumed that all agents can call all other agents.  
The optimal schedule for a 'ring' takes  $2n - 3$  calls.

Example 5 nodes: 6 calls is not possible but 7 calls is possible.

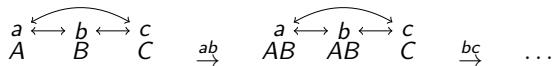
But if the agents can also exchange neighbours (numbers) 6 calls is possible again.



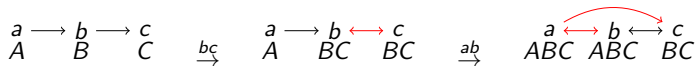
# Dynamic Gossip — Learn New Secrets and Neighbours

*Agents exchange all secrets and all numbers they know.*

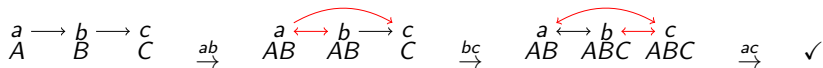
On fully connected graphs there is no difference.



On partially connected graphs LNS deadlock is possible. (After  $bc; ab$  agent  $c$  cannot call agent  $a$ , because  $c$  does not have  $a$ 's number.)



But sometimes LNS deadlock can be avoided.  
(After  $ab; bc$  agent  $a$  calls agent  $c$ .)



When exactly? And how about other protocols?

## Dynamic Gossip — Characterization of success

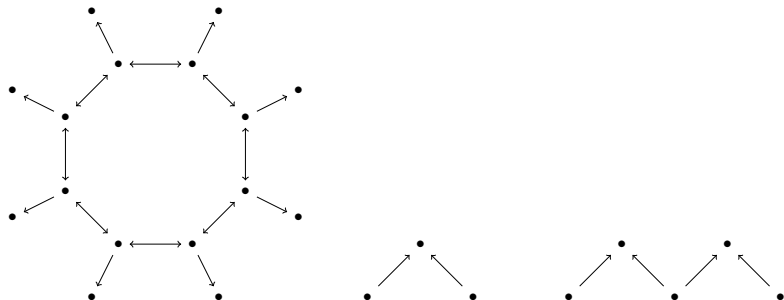
- ▶ A graph is weakly/strongly connected if there is an undirected/directed path between any two nodes.
- ▶ For gossip graphs: properties of the neighbour relation.
- ▶ No protocol is successful on a disconnected gossip graph.
- ▶ All presented protocols except LNS are (strongly or fairly) successful on weakly connected gossip graphs.

*Relevant properties to show these results (let  $G = (A, N, S)$ ):*

- ▶  $S \subseteq N$
- ▶  $S^\sigma \subseteq N^\sigma$
- ▶  $S^\sigma \circ N \subseteq N^\sigma$  ( $S^{\sigma.ab} \subseteq N^\sigma$ )
- ▶ stable  $\tau \subseteq \sigma^\infty$  satisfy  $S_x^\tau = S_y^\tau$ .
- ▶ *[in TOK/SPI every agent is a neighbour of a token holder]*

# Dynamic Gossip — Characterization of LNS success

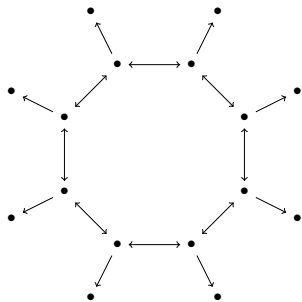
- ▶ If graph **not weakly connected**, unsuccessful. *Worse:*  
If graph a **bush** or **double bush**, unsuccessful.
- ▶ If graph **strongly connected**, strongly successful. *Better:*  
If graph a **sun**, strongly successful.



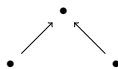
[vD, van Eijck, Pardo, Ramezani, Schwarzenruber. *Dynamic Gossip*. Bulletin of the Iranian Mathematical Society, 2019]

# Dynamic Gossip — Characterization of LNS success

- ▶ A **sun** is a strongly connected graph to which may be linked maximal nodes.
- ▶ A **bush** is a tree with branching factor in the root at least 2.
- ▶ A **double bush** consists of two bushes joined in a leaf linked to their roots.



sun



bush



double bush

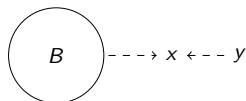
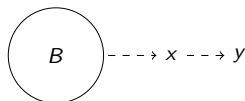
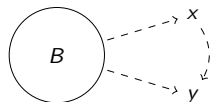
# Dynamic Gossip — Characterization for LNS

- ▶ *LNS is (strongly) successful on a gossip graph iff the gossip graph is a sun*

Lemmas to obtain this result:

- ▶ if  $\sigma$  is LNS-maximal, then  $S^\sigma = N^\sigma$ .  
(otherwise there are  $a, b$  with  $N^\sigma ab$  but  $\neg S^\sigma ab$ , so  $\sigma.ab$  OK)
- ▶ if  $\sigma$  is LNS-maximal, then  $S^\sigma \circ N^* = S^\sigma$   
(sketch:  $S^\sigma \circ N \subseteq N^\sigma = S^\sigma$ , iterate this ...).

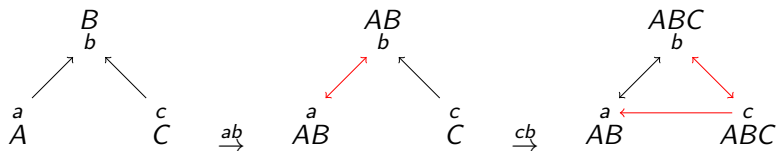
In all cases below there are LNS-maximal  $\sigma$  where  $y$  not expert:



# Dynamic Gossip — Characterization for LNS

- ▶ *LNS is unsuccessful on a weakly connected gossip graph iff the weakly connected gossip graph is a bush or a double bush*

For example, LNS is unsuccessful on this bush:



The (rather) hard part of the proof is to construct a successful call sequence on an arbitrary weakly connected graph that is not a bush or a double bush.

[vD, [van Eijck](#), [Pardo](#), [Ramezani](#), Schwarzentruher. *Dynamic Gossip*. Bulletin of the Iranian Mathematical Society, 2019]



# Dynamic Gossip — Characterization for LNS

- ▶ *LNS is unsuccessful on a weakly connected gossip graph iff*

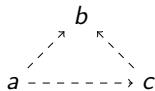
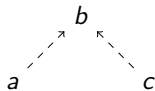
*the weakly connected gossip graph is a bush or a double bush*

( $\Leftarrow$ ) [the easy part] Given bush  $G = (A, N, S)$  with root  $r$ , and LNS-permitted  $\sigma$ . Then:

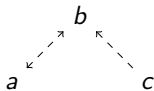
1.  $G|N_x^\sigma$  is a tree.
2.  $N_x^\sigma \setminus S_x^\sigma = \begin{cases} \text{root of } G|N_x^\sigma & \text{if } \neg S^\sigma xr \\ \emptyset & \text{otherwise} \end{cases}$

Sketch: Once the first child of the root calls the root, the subtrees generated by all other children are blocked (not LNS-permitted).

( $\Leftarrow$ ) [the hard part, by example only, see BIMS 2019 for full proof]



$ab; ac; bc$



$cb; ab; ca$

# Dynamic Gossip — Characterization for LNS

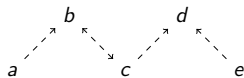
Adding one edge to a double bush makes LNS successful ...



unsuccessful



$cb; ab; cd; ed; ad; bd; ca; ea$



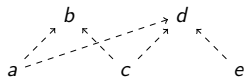
$ab; cd; ed; db; cb; ac; eb$



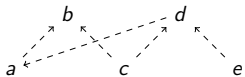
$ab; cd; ed; db; cb; ac; eb$



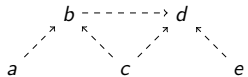
$ab; cd; ed; da; ca; eb$



$ab; cd; ed; ad; bd; cb; eb$



$ab; cd; ed; da; ca; eb$



$ab; cd; ed; ad; bd; cb; eb$



$ab; ae; be; cb; bd; ad; cd; ed$

All for today!