# Functional Programming for Logicians - Lecture 1 Functions, Lists, Types

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module L1 where

Introduction

## Who is who

### Who is who

#### You

- ▶ a wide range of programming experiences: nothing, Java, Python, Rust, Agda, Lean, Prolog, Lisp, C, C++, C#, Haskell, Dart, Vala, Kotlin, Mathematica, ...
- interests: Category Theory, Cognition, Dynamic Epistemic Logic, Inquisitive Semantics, Proof Theory, Recursion Theory, Truth Makers, . . .

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Malvin

- 2012–2014 MoL
- 2014–2018 PhD at ILLC
- 2018–2021 PostDoc in Groningen
- 2021– assistant prof at ILLC

## **Functional Programming**

- the main operation is function application
- describe what, not how it should be computed
- a program is a list of definitions of functions

### Haskell



- Iambda calculus meets category theory
- **typed**: every expression has a type fixed at compile time
- lazy: only compute what and when it is needed
- **pure**: functions have no side-effects
  - ► same input → same output

# Why?

Why?



(Simon Peyton-Jones: *Escape from the ivory tower: the Haskell journey*)

# Let's go

### Calculating

We work in *ghci* for now, the *interactive* compiler.

λ> 7 + 8 \* 9
79
λ> (7 + 8) \* 9
135
λ> sum [1,6,10]
17

Create a file example.hs which contains this:

square x = x \* x

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```
λ> square 9
81
λ> square 10
100
```

Create a file example.hs which contains this:

```
square x = x * x
```

Now we can run ghci example.hs and use this function!

```
λ> square 9
81
λ> square 10
100
```

 $\Rightarrow$  How can we define double, cube and plus?



Our first Type (Error)

with actual type '[Char]'

```
Our first Type (Error)
```



The definition of square we were actually using is this:

square :: Integer -> Integer
square x = x \* x

```
Our first Type (Error)
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```
square :: Integer -> Integer
square x = x * x
```

We read the :: double colon as "has the type"

In Haskell everything has a type!

```
Our first Type (Error)
```





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square :: Integer -> Integer
square x = x \* x

We read the :: double colon as "has the type"

In Haskell everything has a type!

 $\Rightarrow$  What are the types of 10, "10", +, \* and +5?



#### Lists

```
myList :: [Integer]
myList = [1,23,42,111,1988,10,29]
longList :: [Integer]
longList = [1..100]
\lambda> length myList
7
\lambda> length longList
100
\lambda> 1:3:myList
[1,3,1,23,42,111,1988,10,29]
\lambda> myList ++ [5,7] ++ myList
[1, 23, 42, 111, 1988, 10, 29, 5, 7, 1, 23, 42, 111, 1988, 10, 29]
```

#### mapping over lists

```
λ> map square myList
[1,529,1764,12321,3952144,100,841]
```

```
\lambda> map square [1..4]
[1,4,9,16]
```

```
λ> map (*5) [1,2,3,5]
[5,10,15,25]
```

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```
\Rightarrow What does map do?
```



 $\Rightarrow$  What is the type of map? Here? In general?

### mapping over lists

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\lambda> map square [1..4]
[1, 4, 9, 16]
\lambda> map (*5) [1,2,3,5]
[5, 10, 15, 25]
\Rightarrow What does map do?
\Rightarrow What is the type of map? Here? In general?
```

How can we define map?



Type Variables and Inference

```
wordList :: [String]
wordList = ["beyonce", "metallica", "k3", "anathema"]
\Rightarrow Why does map square wordList give an error?
Hint: Look at the error generated by this:
\lambda> import Data.Char
\lambda> :t toUpper
toUpper :: Char -> Char
```

 $\lambda$ > map toUpper wordList

. . .

## The List Monster



#### $\Rightarrow$ Define these four functions, start with the type!

picture from http://learnyouahaskell.com/starting-out/#an-intro-to-lists

Strings are lists of characters

In fact we have:

```
type String = [Char]
```

Example:

λ> "barbara" == ['b','a','r','b','a','r','a'] True Strings are lists of characters

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Example:

λ> "barbara" == ['b','a','r','b','a','r','a'] True

Note the difference between ' and ":

λ> :t 'a'
'a' :: Char
λ> :t "a"
"a" :: [Char]

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λ> :t 'a'
'a' :: Char
λ> :t "a"
"a" :: [Char]

 $\Rightarrow$  Why does 'ab' not make sense?



Mapping and Sorting Strings

```
swab :: Char -> Char
swab 'a' = 'b'
swab 'b' = 'a'
swab c = c
```

# Mapping and Sorting Strings

```
swab :: Char -> Char
swab 'a' = 'b'
swab 'b' = 'a'
swab c = c
\lambda> map swab "abba"
"baab"
\lambda> map swab "barbara"
"abrabrb"
```

# Mapping and Sorting Strings

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swab :: Char -> Char
swab 'a' = 'b'
swab b' = a'
swab c = c
\lambda> map swab "abba"
"baab"
\lambda> map swab "barbara"
"abrabrb"
\lambda> import Data.List
\lambda> sort "hello"
"ehllo"
\lambda> sort "barbara"
"aaabbrr"
```

# Infinite Lazy Lists

What happens here?

```
naturals :: [Integer]
naturals = [1..]
```

What happens if I evaluate naturals in ghci now?

Hint: Maybe I shouldn't



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But we can ask for finite parts of it, lazily!

```
\lambda> take 11 naturals
[1,2,3,4,5,6,7,8,9,10,11]
\lambda> map square (take 11 naturals)
[1,4,9,16,25,36,49,64,81,100,121]
```

## Infinite Lazy Lists

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naturals = [1..]
```

What happens if I evaluate naturals in ghci now?

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[1,2,3,4,5,6,7,8,9,10,11]
\lambda> map square (take 11 naturals)
[1,4,9,16,25,36,49,64,81,100,121]
```

```
\lambda> take 11 (map square naturals) -- not strict!
[1,4,9,16,25,36,49,64,81,100,121]
```

```
\Rightarrow exercise: Give a definition of take.
```

#### Recursion

```
sentence :: String
sentence = "Sentences can go " ++ onAndOn where
onAndOn = "on and " ++ onAndOn
```

Try this out with take 65 sentence in ghci.
# Туре Нуре

- ▶ Integer
- Int
- ▶ [a]
- Char
- String = [Char]

# Туре Нуре

- Integer
- Int
- ▶ [a]
- Char
- String = [Char]

Tuples (aka products):

(a,b)
(a,b,[c])

Sum types:

Either a bMaybe a

## Tuples

```
malvin, jana :: (String, Integer)
malvin = ("Malvin", 1988)
jana = ("Jana", 1993)
```

### **Tuples**

```
malvin, jana :: (String,Integer)
malvin = ("Malvin",1988)
jana = ("Jana",1993)
```

Can you guess what the following functions do?

```
fst :: (a,b) -> a
snd :: (a,b) -> b
Data.Tuple.swap :: (a,b) -> (b,a)
```

# **Tuples**

```
malvin, jana :: (String, Integer)
malvin = ("Malvin", 1988)
jana = ("Jana", 1993)
```

Can you guess what the following functions do?

#### Lambdas

We write  $\setminus$  for  $\lambda$  to define an anonymous function:

```
λ> (\y -> y + 10) 100
110
λ> map (\x -> x + 10) [5..15]
[15,16,17,18,19,20,21,22,23,24,25]
```

#### Lambdas

We write  $\setminus$  for  $\lambda$  to define an anonymous function:

λ> (\y -> y + 10) 100
110
λ> map (\x -> x + 10) [5..15]
[15,16,17,18,19,20,21,22,23,24,25]

 $\Rightarrow$  How can we define fst, snd and swap with lambdas?

Function application and composition

```
people :: [(String,Integer)]
people = [jana,malvin]

λ> map (length . fst) people
[4,6]
```

```
\lambda> concat $ map fst people
"JanaMalvin"
\lambda> sum $ map snd people
3981
```

Function application and composition

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```



- What do . and \$ do?
- Why is \$ still useful?
- Why should we call (length . fst) "point-free"?

# List Comprehension

We can also build new lists using this notation:

threefolds :: [Integer]
threefolds = [ n | n <- [0..], mod n 3 == 0 ]</pre>

The notation is close to *set* comprehension:

 $\{n \in \mathbb{N} \mid n \equiv 0 \mod 3\}$ 

# List Comprehension

We can also build new lists using this notation:

threefolds :: [Integer] threefolds = [  $n \mid n \leq [0..]$ , mod  $n \mid 3 \equiv 0$  ]

The notation is close to *set* comprehension:

$$\{n \in \mathbb{N} \mid n \equiv 0 \mod 3\}$$

An equivalent way to define the above:

filter ( $n \rightarrow mod n 3 == 0$ ) [0..]

### Even more Lists

These are all the same:

```
[1..10]
[1,2,3,4,5,6,7,8,9,10]
1:2:3:4:5:6:7:8:9:10:[]
1:2:3:4:5:6:[7..10]
[ x | x <- [1..100], x <= 10 ]
takeWhile (< 11) [1..]</pre>
```

## Even more Lists

These are all the same:

```
[1..10]
[1,2,3,4,5,6,7,8,9,10]
1:2:3:4:5:6:7:8:9:10:[]
1:2:3:4:5:6:[7..10]
[ x | x <- [1..100], x <= 10 ]
takeWhile (< 11) [1..]</pre>
```

But what about this one?

```
filter (< 11) [1..]
```

Is it the same value? Is it the same program?



# Guards



# Guards



... we usually prefer guards like this:

 $\Rightarrow$  What is the type of otherwise and what does it do?



#### How to make a type

type defines types that are *just abbreviations*:

```
type Person = (String,Integer)
type Group = [Person]
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To create actually new types we use data:

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data MyEither a b = MyLeft a | MyRight b
data MyMaybe a = MyNothing | MyJust a

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data Animal = Cat | Horse | Koala
data MyEither a b = MyLeft a | MyRight b
data MyMaybe a = MyNothing | MyJust a
```

This defines a new type and constructors at the same time!

# Pattern matching

Each data type can be matched by patterns:

- Bool: True, False, b
- Lists: [], (x:xs), (x:y:rest), ...
- Strings: 'h':'e':[], "hello", ...
- Tuples: (x,y)
- Numbers: 0, 1, 2, 3, 42, ...
- Maybe a: (Just x), Nothing
- Either a b: Left x, Right y
- anything: x, mySuperLongVarName, \_

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Patterns can occur in two places:

as arguments of functions:

isEmpty :: [a] -> Bool
isEmpty [] = True
isEmpty (\_:\_) = False

▶ in case ... of ... -> ... constructs.

Logic in Haskell

# **Propositional Logic**

Propositional Logic formulas are defined by:  $\varphi ::= p_n \mid \neg \varphi \mid \varphi \land \varphi$ In Haskell:

data Form = P Int | Neg Form | Conj Form Form

# **Propositional Logic**

Propositional Logic formulas are defined by:  $\varphi ::= p_n | \neg \varphi | \varphi \land \varphi$ In Haskell:

data Form = P Int | Neg Form | Conj Form Form

Given an assignment  $v \colon P \to \{\top, \bot\}$ , we define:

$$v \vDash p_i : \iff v(p_i) v \vDash \neg \varphi : \iff \text{not } v \vDash \varphi v \vDash \varphi \land \psi : \iff v \vDash \varphi \text{ and } v \vDash \psi$$

type Assignment = Int -> Bool

satisfies :: Assignment -> Form -> Bool
satisfies v (P k) = v k
satisfies v (Neg f) = not (satisfies v f)
satisfies v (Conj f g) = satisfies v f && satisfies v g

### Examples

We define an assignment:

world :: Assignment
world 0 = True
world 1 = False
world 2 = True
world \_ = False

### Examples

We define an assignment:

```
world :: Assignment
world 0 = True
world 1 = False
world 2 = True
world _ = False
\lambda> satisfies world (Neg . Neg $ P 2)
True
\lambda> satisfies world (Conj (Neg $ P 1) (P 0))
True
```

Actually, you want this:

```
data Form = P Int | Neg Form | Conj Form Form
  deriving (Eq,Ord,Show)
```

Eq, Ord and Show are type classes, a topic for tomorrow.

**Practical Stuff** 

# Abbreviation Mania

- GHC is the Glasgow Haskell Compiler
- GHCi is the interactive interface of GHC
- stack is a build tool to simplify your life
- cabal is another tool and a package format
- Hackage is a public database of Haskell libraries
- Stackage provides stable snapshots, called *resolvers*.
- ► *VS Code* is a common and beginner-friendly editor.

# Organization

Course website: https://malv.in/2022/funcproglog

Lectures on Zoom

Exercise Sessions on gather.town

 $\Rightarrow$  Links are in the first email!

# Organization

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- Lectures on Zoom
- Exercise Sessions on gather.town
- $\Rightarrow$  Links are in the first email!
- Other useful tips for all-online:
  - Try out screen-sharing in gather.town
  - ▶ If you want to work live in pairs, use a collaborative editor!
    - EtherPad (plain text) on gather.town tables!
    - "Live Share" extension for VS Code
    - https://replit.com/

You can download the lecture and exercises as .1hs (or .hs) files.

This stands for "Literate Haskell" and is a way to combine programs and documentation or longer comments in one file.

In .1hs files the actual Haskell code has to be

- indented with with > (Markdown style) or
- between \begin{code} ... \end{code} (MTEXstyle)

#### How to start

- 0. See instructions in first email to install Haskell and VS Code.
- 1. Download E1.1hs and open a terminal where you saved it.
- 2. Run ghci E1.lhs (or stack ghci E1.lhs).
- 3. Edit the file in VS code.
- 4. Reload with :r and read carefully what GHC tells you.
- 5. Try out all the things!
- 6. Go to 3.



#### See you again at 13:00 in gather.town.